

APPLICATION FOR A UNITED STATES PATENT

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5 **Title:** **Apparatus and Method for a Digital Delay Locked Loop**

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FIELD OF THE INVENTION

This present invention relates to digital circuits. More specifically, it relates to delay locked loop design.

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BACKGROUND OF THE INVENTION

Accurate high-speed clock signals are often used for transmitting and receiving data in high-speed circuits, such as high-speed memory devices or high speed bus channels. However, high speed clock signals often become distorted during transmission and reception. Consequently, high speed circuits, such as dynamic random access memory (DRAM) devices, often recover an externally provided high-speed clock signal by locking an accurate internally generated clock signal to the distorted externally provided clock signal.

A delayed locked loop (DLL) is typically used to delay the internally generated clock signal in order to match the phase of the internally generated clock signal to the phase of some reference clock signal. Typically, a phase-detection circuit in the DLL compares the phase of the internal clock signal to the reference signal and a control-logic block that is coupled to the output of the phase-detection circuit is used to increase or decrease a delay produced by a chain of delay elements used to delay the internal clock signal. U.S. Patent Nos. 5,945,862 and 6,125,157 to Donnelly et al. represent two

approaches to locking an internal clock signal to an external clock signal using delay elements.

It is desirable to provide for a fully digital DLL circuit that can be fabricated using standard digital design techniques.

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SUMMARY OF THE INVENTION

In accordance with a first aspect of the present invention, a circuit for selectively delaying a reference clock signal is provided. The circuit includes a phase splitter coupled to a first delay line and a second delay line. Each delay line includes a set of
5 output taps and has a multiplexor that is selectively coupled to the output taps. An output of each multiplexor is coupled to a latch. The latch provides an output clock signal, which may be fed back to a phase detector, where the phase detector also receives an input clock signal. A controller is coupled to the phase detector and, depending upon the output of the phase detector, provides control signals to each multiplexor.

10 In a preferred embodiment, the circuit is implemented using standard digital design techniques.

In accordance with a second aspect of the present invention, a method for recovering a clock signal from an input clock signal is provided. The method includes converting a reference clock signal into an in-phase reference and a complementary
15 reference, and selectively delaying the in-phase reference and the complementary reference. The method further includes generating an output clock signal from the selected in-phase reference and the selected complementary reference.

In accordance with the preferred embodiment, the selective delay of the in-phase reference and the complementary reference is based upon a difference signal determined
20 by comparing an input clock signal with a feedback clock signal, where the feedback clock signal is related to the output clock signal.

In accordance with another embodiment, the method is implemented using digital design techniques.

This summary is not intended to be all-inclusive, but rather illustrative. These and other aspects of the present invention, and its various embodiments, are described in
5 greater detail below.

BRIEF DESCRIPTION OF THE DRAWINGS

Preferred embodiments of the present inventions are described with reference to the following drawings, wherein:

FIG. 1 is a functional block diagram of one embodiment of a DLL circuit
5 according to the present invention;

FIG. 2 is a timing diagram illustrating an example of the function of the DLL circuit of FIG. 1;

FIG. 3 is a control flow diagram illustrating an embodiment of the function performed by the controller of FIG. 1;

10 FIG. 4 is a logic circuit diagram illustrating an embodiment of a circuit for generating a quadrature clock according to one aspect of the present invention;

FIG. 5 is a control flow diagram illustrating an embodiment of the function performed by a controller to generate control signals to obtain the quadrature clock signal from the circuit of FIG. 4;

15 FIG. 6 is a timing diagram illustrating an example of the function of the quadrature clock circuit of FIG. 4;

FIG. 7 is a functional block diagram of another embodiment of a DLL circuit according to the present invention;

20 FIG. 8 is a functional block diagram illustrating an example of a delay element suitable for use in the delay line circuits of FIG. 7;

FIG. 9 is a control flow diagram illustrating an embodiment of the function performed by the controller of FIG. 7;

FIG. 10 is a functional block diagram illustrating an example of a circuit for detecting a cycle boundary in the delay line circuits of FIG. 7;

5 FIG. 11 is a transistor circuit diagram which illustrates an example of a prior-art tri-state inverter circuit configured to be enabled by a differential enable signal that may cause a multiplexor circuit which uses it to experience a data dependent response;

FIG. 12 illustrates a circuit that addresses the problem of data dependent response exhibited by the circuit of FIG. 11;

10 FIG. 13 is a functional block diagram illustrating an example of a two level multiplexor design suitable for use with the circuits of FIG. 1 and FIG. 7; and

FIG. 14 is a logic circuit diagram illustrating an embodiment of a RS flip-flop suitable for use in the described embodiments.

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DETAILED DESCRIPTION OF PREFERRED EMBODIMENTS

In one preferred embodiment, a reference clock having a correct duty cycle, e.g. 50%, is split into an in-phase clock reference signal and a complementary reference clock signal, each of which is successively delayed using a pair of delay lines. A set of output taps on the delay line allows the selection of and variation in the length of the delay. A pair of multiplexors (MUXes) is used to select one output tap from each of the delay lines. In accordance with one preferred embodiment, a rising edge of a signal from the output tap selected from the delay line used to delay the in-phase reference signal is used to produce a rising edge in an output clock signal. Then, a rising edge of a signal from the output tap selected from the delay line used to delay the complementary reference signal is used to produce a falling edge in the output clock signal. A feedback clock signal is compared to the input clock signal in order to obtain a difference signal that reflects the phase relationship between the feedback clock signal and the input clock signal. A control circuit monitors the difference signal and generates a control signal that drives the selection performed by the MUXes such that a rising edge of the feedback clock is aligned or lock-in to a rising edge of the input clock signal.

In another aspect of the preferred embodiment, a power control signal controls delay elements within the delay lines and the control circuit is configured to generate the power control signal based upon the control-signal provided to the MUXes such that unused elements in the delay lines are disabled during operation to reduce power consumption.

In still another aspect of the present invention, four additional MUXes are used to select output taps from each of the delay lines in order to produce a quadrature clock signal. The additional outputs taps are selected such that they are offset from the previously selected output taps by an offset value, where one offset output tap is
5 advanced by the offset value and the other offset output tap is delayed by the offset value. The additional output taps are interpolated in order to generate the quadrature clock. A quadrature control circuit varies the offset value and monitors the relationship between edges in the signals from the offset output taps in order to drive the additional pairs of MUXes. The quadrature control circuit selects an offset value that results in the edges
10 from the offset output taps being generally aligned.

FIG. 1 is a functional block diagram of one embodiment of a delay locked loop ("DLL") circuit according to the present invention. A reference clock signal REFCLK is received by a phase splitter 102, which splits the REFCLK into a positive phase reference signal RP (also referred to as an in-phase reference signal) and a negative phase reference
15 signal RN that is 180° out of phase with RP (also referred to as the complementary reference signal). The positive phase reference RP is input to a delay line 110, which has N output taps that are input to a multiplexor (MUX) 112. The MUX 112 also receives a CONTROL signal that selects one of the N inputs from the delay line 110 to be output from the MUX 112 as a ZUP signal to set an S input of a Set/Reset latch 130. The delay
20 line 110 successively delays the reference signal RP such that each one of the N output taps represents RP delayed by a multiple of a predetermined delay interval. Thus,

selecting one of the N output taps from the delay line 110 using the MUX 112 permits the RP signal to be delayed by from 1 to N delay intervals before being input to the S terminal of the latch 130.

Likewise, the delay line 120 delays the negative phase reference signal RN by from 1 to N delay intervals, as selected by the MUX 122. The delay line 120 receives RN and has N output taps that are input to the MUX 122. The MUX 122 also receives the CONTROL signal, which determines which of the N output taps of the delay line 120 to output as a ZDOWN signal to a reset input R of the latch 130. The MUX 122 is configured to respond to the CONTROL signal in the same fashion to the response of the MUX 112. In other words, the delay line 120 and the MUX 122 are configured relative to the delay line 110 and the MUX 112 such that the ZUP and ZDOWN signals are complementary to one another in generating signals responsive to the CONTROL signal output by controller 140. The ZUP and ZDOWN signals control the up/down edges of the output signal. The output signal drives some collection of circuitry to produce a signal FB_CLK that is aligned in zero degrees in phase with the input clock signal. The resulting output clock signal CLKOUT generated by latch 130 is a clock signal whose edges are determined by the CONTROL signal and which has a 50% duty cycle.

A phase detector circuit 142 receives a feedback clock signal FB_CLK derived from the CLKOUT signal and also receives an input clock signal INPUT CLK that is the signal to which the circuit 100 is to lock onto. The CLK_OUT may move in time in a

manner that the resulting phase relationship of the FB_CLK relative to the INPUT CLK is well managed.

The phase detector 142 compares the FB_CLK and the INPUT CLK and outputs a difference signal DIFF that indicates whether the phase of FB_CLK leads or lags the phase of INPUT CLK. The DIFF signal is received by the controller 140, which adjusts the CONTROL signal accordingly. As one example for the control circuit 140, the circuit may consist of a counter, which is incremented or decremented by the DIFF signal in order to adjust the CONTROL signal. It should be understood that someone skilled in art of PLL or DLL design will recognize that the circuit which provides the control signal to the VCO or delay-line of a DLL can be implemented in many fashions, and other circuits of varying complexity could also be used.

FIG. 2 is a timing diagram illustrating an example of the function of the DLL circuit of FIG. 1. In FIG. 2, the reference clock signal REFCLK is shown along with the reference clock signals RP and RN output from the phase splitter 102. Note that REFCLK has a 50% duty cycle. Also note that RP and RN are complementary to one another, e.g. 180° out-of-phase with one another and that a delay is introduced to RP and RN relative to REFCLK by the phase splitter 102.

FIG. 2 also includes an example of an external input clock signal INPUT CLK. Note that, in this example, the INPUT CLK signal waveform is distorted due to the effects of transmission and, consequently, the received INPUT CLK waveform does not have a 50% duty cycle. In one embodiment, phase detector 142 compares the phase of

rising edge 150 of INPUT CLK to the FB CLK, which is derived from CLK OUT. Controller 140 receives the DIFF signal and adjusts the value of the CONTROL signal in order to select a delay tap DP(n) from delay line 110 that aligns the rising edge of FB CLK with the rising edge 150 of INPUT CLK. In this example, DP(n) is selected such
5 that rising edge 152 of DP(n) is aligned with rising edge 150 of INPUT CLK. The rising edge 152 of DP(n), in turn, drives the set input of the latch 130 triggering the rising edge 154 of the CLK OUT signal.

The CONTROL signal will also select DN(n), from the delay line 120, for output by the MUX 122 to the reset terminal of the latch 130. DP(n) and DN(n) are generally
10 complementary to one another, but are typically distorted by their respective delay lines 110 and 120 such that neither DP(n) nor DN(n) will have a correct 50% duty cycle. However, the distortion will effect the rising edges of DP(n) and DN(n) in the same manner. Therefore, while the duty cycles of DP(n) and DN(n) are distorted, the relationship between the rising edge 152 of DP(n) and the rising edge 160 of DN(n) is
15 reliable. Thus, in the present invention, rising edge 160 of DN(n) is used to drive the reset input of the latch 130 triggering a falling edge 162 of the CLK OUT. The cycle then continues as demonstrated by next subsequent a rising edge 170 of DP(n) that triggers a rising edge 172 of the CLK OUT. The result is that the circuit 100 of FIG. 1 may produce a CLK OUT signal having a 50% duty cycle.

20 Note that the timing diagram of FIG. 2 has been simplified to demonstrate the relationship between edges in the various waveforms. The actual waveforms are subject

to further delays introduced, for example, by the MUXes 112 and 122, the latch 130, as well as other circuitry that is intermediate to the CLK OUT signal and the FB CLK signal.

As noted above, the controller 140 may be, for example, a counter, or composed
5 of complex circuitry, such as a processor, that monitors the output of the phase detector 142 and generates the CONTROL signal. FIG. 3 is a control flow diagram illustrating an embodiment of a process 180 performed by the controller 140 for generating the CONTROL signal. In the process 180, at step 182, an initial value of n is selected and, at step 184, the CONTROL signal for controlling the delay of delay lines 110 and 120 is
10 generated according to the initial value of n selected at step 182.

At step 186, the output of the phase detector 142 is checked to determine if signal lock has been achieved between FB CLK and INPUT CLK, e.g. the rising edge of FB CLK is aligned with the rising edge of INPUT CLK. If signal lock is achieved, then control flow branches at step 190 to step 194, where the process waits for a short time
15 interval before branching back to step 186 to check the output of the phase detector 142 again.

If signal lock is not achieved, then control flow branches at step 190 to step 192 where the value of n is adjusted according to the phase detector output. For example, if the DIFF signal is negative, then n is decremented. If the DIFF signal is positive, then n
20 may be incremented. Process 180 then branches back to step 184, where the CONTROL signal is generated according to the adjusted value of n .

Note that other approaches are also possible. For example, the signal lock branch step 190 and wait step 192 may be omitted so that the value of n is always adjusted. In this case, when signal lock is achieved, the value of the DIFF signal will toggle between positive and negative on each loop through the process. Control for DLLs is well understood in the art and one of ordinary skill will readily appreciate that the circuit of FIG. 1 may be adapted to work with a variety of approaches to DLL design.

In another embodiment, a quadrature clock signal may also be generated. A quadrature clock is a clock that is out-of-phase by 90° or 270° from a main clock signal. Figure 4 illustrates an embodiment of a quadrature clock generation circuit 200 that may be used in combination with the circuit of FIG. 1. The quadrature clock generation circuit 200 includes a MUX 210 that receives N output tap lines of the delayed signal DP from the delay line 110, and is controlled by a control signal XDPCTRL1, which selects one of the taps DP($n+X$) for output to a first input of an interpolator circuit 262. The interpolator circuit 262 illustrated in FIG. 4 is composed of three inverters 212, 220 and 232. Interpolator circuits are generally understood to those of ordinary skill in the art, and other interpolator designs may be adapted for use in the present invention. The circuit 200 further includes a second MUX 230, which receives the N output tap lines of the delayed signal DN from the delay line 120 and a control signal XDNCNTRL1, which selects one of the taps DN($n-X$) for output to a second input of the interpolator circuit 262.

The interpolator circuit 262 interpolates between a rising edge of the DP signal and the rising edge of the DN signal, i.e., it takes an average of the 0-degree signal and the 180-degree signal to generate a 90-degree QUP signal. To do that, the interpolator 262 uses a late version of the DP signal and an early version of the DN signal, i.e., the signals DP(n+X) and DN(n-X), respectively.

The inverter 212 of the interpolator circuit 262 receives the DP(n+X) signal, and the inverter 232 receives the DN(n-X) signal. The outputs of both inverters are coupled to the input inverter 220, which possesses a parasitic capacitance. The inverters 212 and 232 are relatively small devices relative to the inverter 220. For instance, the inverter 200 may be twice the size of the inverters 212 and 232. The inverters 212 and 232 will charge the parasitic capacitance of the input of the inverter 220 until a threshold voltage for the inverter 220 is reached causing its output to change state in order to generate a QUP signal to set an S input of an RS latch 260. When the signals driving the inverters 212 and 232 are separated in time by less than the time it takes for the inverter 212 to fully transition the output of the inverter 220, then the inverter 232 contributes to the voltage transition at the output of the inverter 220. Thus, the interpolator circuit 262 will interpolate between a late version of DP signal and an early version of the DN signal in order to obtain a 90-degree QUP signal.

Further, as illustrated in FIG. 4, the circuit 200 includes a second pair of MUXs 244 and 246, and a second interpolator circuit 264 composed of inverters 248, 250, and 252 that generate a 270-degree QDOWN signal. The MUX 244 receives the N output tap

lines of the delayed DP signal from the delay line 110, and is controlled by a control signal XDPCNTRL2, which selects one of the taps DP(n-X) for output to the second interpolator circuit 264. Similarly, the MUX 246 receives the N output tap lines of the delayed signal DN from the delay line 120 and control signal XDNCNTRL2, which
5 selects one of the taps DN(x+N) for output to the interpolator circuit 264.

To generate the QDOWN signal, the interpolator circuit 264 takes an average of the 180-degree signal and the 0-degree signal (or 360-degree signal) to generate a 270-degree QDOWN signal. To do that, the interpolator 264 uses an early version of the DP signal and a late version of the DN signal, i.e., the signals DP(n-X) and DN(n+X),
10 respectively. The QDOWN signal is then input to a reset input R of the latch 130, and the latch 130 outputs a quadrature clock Q_CLK_OUT that is out-of-phase by 90 degrees or 270 degrees from the main clock.

The circuit 200 further includes a quadrature control circuit 240 that generates the four control signals XDPCNTRL1, XDNCNTRL1, XDPCNTRL2, and XDNCNTRL2
15 that are input to the MUXes 210, 230, 244, and 246, respectively. The quadrature control circuit 240 receives a TEST signal along with the value n from the CONTROL signal of the control circuit 140 of FIG. 1 and uses the values of these signals to generate the four control signals for the MUXes. The value of the TEST signal qualifies the value of the offset signal X used by the quadrature control circuit 240. As illustrated in FIG. 4, a
20 phase detection circuit, such as a D flip-flop 242, generates the TEST signal. The DP(n+X) signal selected at the MUX 210 drives a data input terminal of the D flip-flop

242, while the $DN(n-X)$ selected at the MUX 230 drives a clock input terminal of the D flip-flop 254. A Q output terminal of the flip-flop 254 generates the TEST signal for input to the quadrature control circuit 240.

According to one embodiment, starting with a small value of X and incrementing it, the flip-flop 242 will change the state from 1 to 0 at some value of X that is used to generate the QUP and QDOWN signals. Similarly, a large value of X could be first selected and then decremented until the transition occurs. Alternatively, the $DN(n-X)$ signal could be input to the D input, and the $DN(n+X)$ could be then input to the clock input of the D flip-flop 242. In such an embodiment, the selected value of X corresponds to a value of X when the flip-flop 242 changes from 0 to 1. Thus, the phase detection circuit detects when the first input is no longer ahead of the second input, at which point the step of incrementing/decrementing the value of X is stopped, and the value of X can be used as an input to the quadrature control circuit 240.

The function of the interpolator of FIG. 4 is further explained below in conjunction with FIG. 6.

In FIG. 4, a change in state of the TEST signal from the Q output terminal of the flip-flop 242 indicates the midpoint between the two signals being interpolated. FIG. 5 is a flow chart illustrating one embodiment of a process 250 performed by the quadrature control circuit 242. The process 250 starts from a minimum value for X, which is initialized at step 252. The quadrature control circuit 240 then generates control signals at step 254 based upon $DP(n+X)$ and $DN(n-X)$, respectively. The functionality of the

control circuit 240 may be integrated with some embodiments of the control circuit 140 of FIG. 1.

The value of the TEST signal is checked at step 256 to determine if the TEST signal has transitioned from logic one to logic zero. If the TEST signal is at logic zero, then the midpoint may have not been reached, and the control flow branches to step 262, where the value of X is incremented. The control flow then branches back to step 254 where a new set of control signals is generated based on the incremented value of X. The value of the TEST signal is again verified at step 256. If the TEST signal has transitioned to the logic zero, then the control flow branches at step 260 to step 264, where the value of X is decremented. The value of X is decremented because the transition to logic zero indicates that the midpoint has been passed and the value of X is decremented, in this embodiment, to obtain a more accurate midpoint value. The value of X is now set for operation, and the QUP and QDOWN signals are generated. As mentioned in the preceding paragraphs, instead of setting the value of X to a minimum value, the value of X may be set to a maximum value and then decremented until the transition is detected.

Interpolation of the 90-degree quadrature clock signal is illustrated in the waveform of the timing diagram of FIG. 6. The edges produced by the DP(n+X) and the DN(n-X) are shown as the value of X is successively incremented (or decremented) until the two edges converge at 90-degrees with the result being a QUP signal.

As noted above, a 270-degree clock signal may alternatively be generated. Whereas the 90-degree QUP signal was derived by interpolating a midpoint that is

delayed from DP(n) and advanced from DN(n), e.g., taking a late DP signal and an early DN signal, e.g., using DP(n+X) and DN(n-X), a 270-degree QDOWN signal can be derived from an early version of DP signal and a late version of the DN signal, e.g., DP(n-X) and DN(n+X). The interpolation of the 270-degree QDOWN signal is also
5 illustrated in the waveforms of the timing diagram of FIG. 6.

In accordance with another embodiment, unused delay elements of the delay lines may be powered down to save power. For example, once the value of n is determined, some or all of the subsequent delay elements in the delay line may be turned off. FIG. 7 is a functional block diagram illustrating another embodiment of a DLL circuit 300
10 according to the present invention. The DLL circuit 300 of FIG. 7 is similar to the DLL circuit 100 of FIG. 1, except that delay lines 310 and 320 are composed of delay elements that may be powered down and a controller circuit 340 is adapted to generate a POWER CONTROL signal for powering down elements of the delay lines 310 and 320. In this embodiment, the controller 340 is obviously more complex than an up/down counter.

15 FIG. 8 illustrates one example of a delay element circuit 350 that can be powered down and is suitable for use in the delay lines 310 and 320 of FIG. 7 and FIG. 1. The circuit 350 is a differential circuit that receives and outputs complementary signals. Inputs IN and INB receive the delayed signal output from a previous delay element in a sequence of the delay line. Outputs NEXT and NEXTB are output, for example, to the
20 next delay element in the sequence of the delay line. OUT and OUTB are output tap signals that are output to the MUXes 112 and 122. The delay for the delay element

circuit 350 is primarily provided by NAND gates 360 and 362. Inverters 356 and 358 primarily help keep the delay pass complementarily. An ENABLE signal is input to each of the NAND gates 360 and 362. When the ENABLE signal is at a logic one, then the input signal at input terminals IN and INB is enabled to pass through the delay element to the next delay element and to the output terminals OUT and OUTB. When the ENABLE signal is at a logic zero, then the input signal at input terminals IN and INB is blocked. The NAND and inverter gates may be selected such that no current is drawn when no data transitions take place, e.g. complementary metal-oxide semiconductor (CMOS) logic.

10 The POWER CONTROL signal output by the control circuit 340 separately drives the ENABLE line of the delay elements in the delay lines 310 and 320. Once the value of n is determined to align the DLL circuit with the INPUT CLK signal, then unused delay elements may be disabled to reduce power consumption.

FIG. 9 is a control flow diagram illustrating an embodiment of a process 370 performed by control circuit 340 for determining the value of n and setting the POWER CONTROL signal. Process 372 is similar to process 180 of FIG. 3. At step 372 (step 182 in FIG. 3) of process 370, the control circuit 340 initializes the value of n , and the POWER CONTROL signal is initialized such that all elements of delay lines 310 and 320 are active. At step 374 (step 184 in FIG. 3), the delay CONTROL signal is generated that drives the MUXes 112 and 122 of the delay circuit 300 to select the output taps from the delay lines 310 and 320. At step 376 (step 186 in FIG. 3), the output from a phase

detector is checked to determine whether the phase of the FBCLK signal lags or leads the phase of the INPUT CLK signal. If the circuit 300 has not locked to the INPUT CLK signal, then control flow branches at step 380 to step 188, where the control circuit 340 adjusts the value of n according to the signal output by the phase detector 142 and the control flow returns to step 384 where the delay CONTROL signal is generated using the new value of n .

If the circuit 300 has locked onto the INPUT CLK signal, then the control flow branches at step 380 to step 382, where the control circuit 340 determines a power control position that may be used to adjust the POWER CONTROL signal to turn off unused delay elements in the delay lines 310 and 320. At step 384, the control circuit 340 adjusts the power control signal to turn off unused elements, and the adjustment is based on the power control position determined at step 382.

According to an alternative embodiment, the N delay stages ahead of the maximum delay stage of interest may be powered-down. The power control circuit monitors the ZLOOP and QLOOP delay line usage and applies a power down signal to positions "MAX+N," where "MAX" is the highest position being used, and N is 1 and greater.

In a DLL circuit having a large delay line, an entire cycle of the reference clock frequency might be contained within the delay line. For instance, each step of a 64-bit delay line might be 50pS, thus, producing a delay of 3.2nS on the entire line. In such an embodiment, if the reference clock has a frequency of 500 MHz, the entire reference

clock cycle might be represented in 2.0nS, or 40 stages, of that delay line, leaving 24 stages of the line unused. Therefore, it might be advantageous to detect a cycle boundary, where a DLL circuit crosses over from a maximum delay value to a minimum delay value in locking onto a signal.

5 FIG. 10 illustrates an example of a circuit 400 for detecting a cycle boundary and generating a CYCLE BOUNDARY signal that may be received by a control circuit adapted, for example, to adjust the CONTROL signal to the MUXes 112 and 122 or the POWER CONTROL signal accordingly. In the circuit 400, a delay line 410 having a series of output taps receives the REFCLK signal. A first delay output tap Delay(0) is
10 output to a MUX 412, and a maximum delay output tap Delay(N-1) is output to a MUX 414. An output of the MUX 412 drives a buffer 416 whose output is received by a phase detector 420. Likewise, an output of the MUX 414 drives a buffer 418 whose output is also received by the phase detector 420. An output of the phase detector 420 is received by a reference control 430 that generates the CYCLE BOUNDARY signal.

15 In one embodiment, the phase detector 420 receives a first output of the delay line 410, a bit position (0), and finds another point on the delay line 410 that matches the bit position (0) in phase. To do that, the phase detector 420 may compare other bit positions, for example sequentially and in order, until the phase detector 420 finds a match. For instance, the phase detector 420 might check a position N to determine if the bit position
20 N is ahead or after the bit position (0), and, if it is not ahead, the phase detector 420 may instruct reference control 430 to increment the bit position to N+1. At some point, if the

reference frequency is fast enough, the phase detector 420 detects a point where the bit position "N" happens before the bit position "0". At that point, the phase detector 420 determines that one-cycle of the reference clock takes N stages of the delay line 410.

According to one embodiment, the cycle boundary number N determined using the cycle boundary detector circuit 400 may be fed back to manage the primary loop circuit 140 described in reference to Figure 1. During a normal operation, the circuit 140 may be arranged to increment its bit position by one bit position every hour in order to track a slowly drifting reference signal. After some time, a loop may approach a last bit position, such as a bit position 64 of a 64-stage delay line, for instance. Instead of approaching the 64th stage of the delay line, at which point the loop becomes unlocked, and the data might be lost while the loop re-acquires, the loop might utilize the number N determined by the cycle boundary detector circuit of Figure 10. Using the value of N, the loop may determine that the bit position "64-N" may be considered as the bit position 64. Therefore, for instance, if $N=40$, instead of moving to position 64 on an "increment" step, the controller may move to position 24. Alternatively, during an increment step, the controller may move to position "64-N+M", where M is a fixed value, such as 2, for instance.

Further, the value of "N" determined using the cycle boundary circuit 400 may also be used by the quadrature loops. For example, if a quadrature loop ("Q-loop") is locked at bit position 60, and the Q-loop controller is adapted to use +/- 10 bit positions to create QUP and QDOWN signals, the quadrature loop might require bit positions 50

and 70. However, in a 64-bit delay line, the Q-loop controller will not be able to use the 70th bit position. However, if $N=40$, the Q-loop may use bit positions 50 and 30.

One problem that can arise in a DLL is that the circuits used to implement the MUX elements may have data dependent response characteristics, e.g. the speed of the response to a logic zero is different than the speed of the response to a logic one. FIG. 11 illustrates an example of an inverter circuit 450 configured to be enabled by a differential enable signal that may experience a data dependent response. In the circuit 450, an enable signal EN, such as a line from the POWER CONTROL signal of the circuit 300 of FIG. 7, drives an NMOS transistor 458 while a complementary enable signal ENB drives a PMOS transistor 452. When EN and ENB are active, then a DATA signal is inverted by a PMOS transistor 454 and an NMOS transistor 456 to drive an output terminal OUT. When EN and ENB are inactive, then the inverter formed by the transistors 454 and 456 is disabled because no current can flow between the power supply rails V_{DD} and V_{SS} . The circuit 450 suffers from data dependent behavior because the amount of capacitance that the non-selected MUX paths contribute to the output node is a function of the data signal driving the non-selected nodes. For example, if a first data input $IN(0)$ is selected, the devices 454 and 456 contribute one value of capacitance to the OUT node if the $DATA(0)$ is low, and another capacitance value if $DATA(0)$ is high. During the operation, the propagation delay thru the MUX element from the selected input to the common output will therefore vary as a function of the data states on the non-selected inputs, thus, resulting in a data dependent response.

FIG. 12 illustrates a circuit 470 that addresses the problem of data dependent response exhibited by the circuit of FIG. 11. In the circuit 470, a first data input IN(0) is input to an inverter 472 which drives one input of a NAND gate 474. An output of the NAND gate 474 drives an input of tri-state inverter 476. The output of the inverter 476 is coupled to a common output circuit node OUT. A second input of the NAND gate 474 is driven by a first enable signal EN(0) derived from the CONTROL signal received by the MUX. The inverter 476 receives its tri-state control from EN(0). Each input to the MUX has a similar circuit, as illustrated by the ellipses leading to gates 482, 484 and 486 for processing Nth data input IN(N) and enable signal EN(N). Via the circuit shown in FIG. 12, all of the non-selected MUX paths contribute the same amount of capacitance load to the output node, regardless of the data values on the non-selected paths. When EN(0) is low, the logic gate 474 blocks any contribution from IN(0) to the logic state at the input of the tri-state inverter 476. If EN(0) is low, the logic gate 474 drives the node at the input of the tri-state inverter 476 to a logic high, removing data dependence described in reference to FIG. 11. Alternatively, the logic gate 474 could be replaced with other logic elements that achieve the same functionality.

FIG. 13 illustrates an example of a two-level 64 input MUX circuit 500 suitable for use in the circuits discussed above. MUX circuit 500 is constructed from a first level of 8-to-1 MUXes 502, 504, 506, 508, 510, 512, 514 and 516 that are driven by most significant three bits of the CONTROL signal CONTROL<5:3>. The output of each of

the first level MUXes is input to a second level 8-to-1 MUX 520 that is driven by the next three least significant bits CONTROL<2:0> and a control bit <3>.

The output of 8-to-1 MUX 520 produces an output MUX_OUT for the 64-input MUX circuit 500. According to an exemplary embodiment, any of the 64 inputs to the MUXes 502-516 can be selected to become the MUX OUT signal. However, the 64 input signals are not just any digital signals. Instead, they correspond to 64 taps of the delay line. Therefore, a signal transition at a tap "0" may occur 50pS before a signal transition at a tap "1", which then transitions 50pS earlier than a signal at a tap "3", and so on. This characteristic of the delay line, where every consecutive signal is delayed with respect to its first neighboring signal, could change if the MUX adds more error than this unit-delay spacing. To minimize the error introduced by the MUX, it is desirable that the change between any two neighboring signals introduced by the MUX is identical.

For instance, the first eight taps of the delay line could be connected to the first MUX 502, and the next eight taps could be connected to the second MUX 504, and so on. In such an embodiment, if the MUX circuit 500 switches between a bit position 4 and a bit position 5, then the only control signal that is changing is the control code on the MUX 502, and the control code on the MUX 520 stays the same since the MUX 520 reads from the same MUX 502 in the two consecutive reads. Therefore, only one MUX is changing in between the bit positions 4 and 5. However, in such a configuration, if the MUX circuit 500 switches between bit positions 8 and 9 (i.e., a crossover point between the MUXes 502 and 504), the control signals not only change on the input MUXes, but

also on the output MUX 520, since the MUX 520 now selects the output from the second MUX 504. Thus, the process of changing from the bit positions 8 and 9 results in two changes rather than one change as in the case of switching from the bit positions 4 and 5. Therefore, such a configuration is not desirable.

5 Instead, according to a preferred embodiment, the delay line taps are wrapped around the eight MUXs 502-516 so that the control code changes only once during any bit position read. The delay line taps are connected to the MUXes in such a way so that the bit position 1 is connected to the first input of the MUX 502, the bit position 2 is connected to the first input of the MUX 504, the bit position 3 is connected to the first
10 input of the MUX 506, and so on until the bit position 8 is connected to the first input of the MUX 516. According to this embodiment, the next bit position, i.e., the bit position 9, is also connected to the MUX 516. Specifically, the bit position 9 is connected to the second input of the MUX 516. Similarly, going up the line of MUXes, the bit positions 10, 11, 12, 13, 14, and 15 are connected to the second inputs of the MUXes 514, 512,
15 510, 508, 506, and 504. Similarly to the bit positions 8 and 9 being connected to the same MUX 516, the bit positions 16 and 17 are connected to the MUX 502. In such a setup, eight inputs to the MUX 502 are bit positions 1, 16, 17, 32, 33, 48, 49, and 64.

Using the MUX setup described in the preceding paragraph, when the circuit 500 switches between bit positions 1 through 8, none of the control signals on the MUXes
20 502-516 are changing, and only the control signal on the output MUX 520 is changing to select a different MUX for each consecutive output. When the circuit crosses from the

bit position 8 to the bit position 9, the control signals to the MUXes 502-516 are changing, but the control signal on the MUX 520 remains the same since the bit position 8 and the bit position 9 are read from the same MUX 516. As the count increases from the bit position 3 to the bit position 16, the MUX 520 moves up the line from the MUX 516 towards the MUX 502 by changing the control signal being input to the MUX 520. In all cases, only one MUX changes in between any two consecutive positions.

As mentioned earlier, the MUXes 502-516 are driven by three most significant bits of the control signal (i.e., CONTROL $\langle 5:3 \rangle$), and the three bits are decoded into one of eight values, and the decoded value selects which input is passed to the MUX's output. For example, if the three bits are "010", then the second input is passed to the output. Similarly, the output MUX 520 is driven by three least significant bits, i.e., CONTROL $\langle 2:0 \rangle$, and a control signal $\langle 3 \rangle$. The extra control signal is used to invert the three least significant bits. If the control signal is low, no inversion is done, and bits "010" in CONTROL $\langle 2:0 \rangle$ indicate that the second input of the MUX 520 drives the MUX_OUT. If the control signal is high, then the bits "010" are inverted to "101", and the fifth input is now selected. In such an embodiment, the control code "000111" selects the eighth delay line position, and the control word "001000" selects the ninth bit position.

FIG. 14 is a logic circuit diagram illustrating an embodiment of an edge-triggered RS latch 550 suitable for use in the embodiments, e.g. as SR latch 130 of FIG. 1 and FIG. 7. The SR latch illustrated in FIG. 14 has a transition delay of a set signal (S) to an output signal (Q) that is equal to a signal transition from a reset signal (R) to the Q signal.

In the SR latch 550, a D flip-flop 552 is clocked by a signal received at the S input, which then clocks a logical one from V_{DD} to output Q of the flip-flop 552. The output Q of the flip-flop 552 is coupled to an input of a NOR gate 570 and further through an inverter 556 to an input of a NAND gate 560. Similarly, a D flip-flop 554 is
5 clocked by a signal received at the R input, which then clocks a logical one from V_{DD} to an output Q of the flip-flop 554. An output Q of the flip-flop 554 is coupled to an input of a NOR gate 572 and through an inverter 558 to an input of a NAND gate 562. The NAND gates 560 and 562 are cross-coupled such that the output of each is input to the other. Likewise, the NOR gates 570 and 572 are cross-coupled such that the output of
10 each is input to the other. The output of the NAND gate 560 drives an inverter 566 that, in turn, drives an output QB. The output of the NAND gate 562 drives an inverter 564 that, in turn, drives an output Q. Similarly, the output of the NOR gate 570 drives an inverter 574 that, in turn, drives the output Q. The output of the NOR gate 572 drives an inverter 576 that, in turn, drives the output QB. The Q output drives an inverted reset
15 input of the D flip-flop 554 and the QB output drives an inverted reset input of a D flip-flop 552. The NAND gates 560 and 562 and the NOR gates 570 and 572 preserve the state of the outputs Q and QB of the RS latch 550, which are determined by the signals input to the S and R inputs.

The wired OR produced by the coupling of the outputs of the inverters 564 and
20 574 and the interpolator such as the wired or produced by the coupling of the outputs of the inverters 566 and 576 accommodates the one gate delay difference introduced by the

inverters 556 and 558. The number of gate delays from each of the S and R inputs to outputs Q and QB is approximately the same, which reduces the possibility of a race condition arising from signals arriving at different points of the circuit at different times.

The foregoing circuit may be readily implemented by those skilled in the art of digital circuit design, and the embodiments described are not intended to be limited to any particular process or fabrication techniques. For instance, a conventional synthesized logic design flow may be used. It should be understood that the present invention is not limited to the circuits, methods and systems described herein. Equivalent circuits will become apparent to those skilled in the art of circuit design upon review of the foregoing.

10 The performance of various types of circuits and systems may be improved through the application of the teachings described herein.

In view of the wide variety of embodiments to which the principles of the present invention can be applied, it should be understood that the illustrated embodiments are examples only, and should not be taken as limiting the scope of the present invention.

15 For example, the circuits may employ different types of transistors and amplifiers to obtain the functions described above and the functions may be achieved with more or fewer elements than those illustrated above. Further, some elements of the embodiments described may be implemented in software, hardware, firmware, or a combination of these approaches. Further, it will be apparent to those of ordinary skill in the art that the

20 teachings of the present invention may be applied to other systems besides RAMBUS based data busses.

The claims should not be read as limited to the described order or elements unless stated to that effect. Therefore, all embodiments that come within the scope and spirit of the following claims and equivalents thereto are claimed as the invention.